

UNIVERSITY OF SWAZILAND
FINAL EXAMINATION, MAY 2005

- Title of Paper : THEORY OF COMPUTATION
- Course number : CS211
- Time allowed : Three (3) hours.
- Instructions : (1) Read all the questions in Section-A and Section-B before you start answering any question.
(2) Answer all questions in Section-A. Answer **any two** questions of section-B. Maximum mark is 100.
(3) Use correct notation and show all your work on the answer script.

This paper should not be opened until the invigilator has granted permission.

SECTION-A

Q1 (marks 6 + 6+ 12). The following languages are given on symbol set $\{0, 1\}$. Assume $w \in \{0, 1\}^*$.

(i). $L1 = \{w, 1 < |w| < 5\}$

(ii). $L2 = \{w, w \text{ starts with one zero and ends with two ones}\}$

(iii). $L3 = \{w, \text{ such that } |w| \bmod 3 \neq 0\}$

The following set of words is given -

$\{\lambda, 0, 1, 0100, 0011, 1100, 00101, 1111101, 11111, 00000, 0010110, 0101011\}$

(a). From the above set, write all the words belonging to $L1$, all the words belonging to $L2$ and all the words belonging to $L3$.

(b). Write three regular expressions representing $L1$, $L2$ and $L3$.

(c). Design three deterministic finite acceptors (dfa's) that accept $L1$, $L2$, and $L3$.

Q2 (marks 6 + 8 + 12). The following non deterministic finite acceptor (nfa) is given :

$M = (\{q0, q1, q2\}, \{0, 1\}, q0, \delta, \{q1\})$, where the transitions are given as :

$$\delta(q0, 0) = q1 ; \quad \delta(q0, \lambda) = q1 ;$$

$$\delta(q1, 0) = q0 ; \quad \delta(q1, 0) = q2 ; \quad \delta(q1, 1) = q1 ;$$

$$\delta(q2, 0) = q2 ; \quad \delta(q2, 1) = q1 .$$

(i). Draw the transition digraph of the nfa.

(ii). Write the corresponding right linear grammar generating $L(M)$

(iii). Convert the above nfa into an equivalent dfa.

SECTION-B

Note: Answer any two questions in this section.

Q3(a) (marks 10). The regular expression of signed real type data values in Pascal is -

$$(\lambda + s) (dd^*tdd^* + dd^*t+ tdd^*) (\lambda + x (\lambda + s) dd^*)$$

where digit, $d \in \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$, sign, $s \in \{+, -\}$, exponent, $x \in \{e, E\}$ and $t = '.'$.

Give one example of each of the possible forms of real data values. Write a linear grammar for the above language.

Q3(b) (marks 15). Find Context Free (CFG) grammars that generate the following languages -

(i). $L(G1) = \{a^n b^m c^{n+m} \text{ where } n, m \geq 0\}$

(ii). $L(G2) = \{a^n b^m c^k, \text{ either } (n = m) \text{ or } (k \geq m) \text{ where } n, m, k \geq 0\}$

write left most derivations for -

$w1 = aabbccccc$ (using $G1$) and $w2 = abbccc$ (using $G2$).

Include production number of your grammar at each step of derivation.

Q4(a) (marks 15). Design a non deterministic pushdown automaton (npda) to recognize the language -

$$L = \{ w \in \{a, b\}^* , n_a(w) = n_b(w) , w \text{ always starts with an } a \}$$

Describe the functional steps of your npda. Write instantaneous descriptions for $w = aabbab$.

Q4(b) (marks 10). Design a non deterministic pushdown automaton (npda) to recognize the language generated by the grammar in Greibach Normal Form-

$$G = (\{S, A, B, C\}, \{a, b, c\}, S, P)$$

where the set of productions P is

$$\left\{ \begin{array}{l} S \longrightarrow aA \\ A \longrightarrow bB \mid aABC \mid a \\ B \longrightarrow b \\ C \longrightarrow c \end{array} \right\}$$

Write instantaneous descriptions for $w = aaabc$.

Q5 (marks 15 + 5 + 5). Write the functional steps of the design of a Turing machine to compute -

$$F(x) = x \text{ div } 3$$

Assume x to be a non zero positive integer in unary representation. Also write the design and instantaneous descriptions using the values of x as 1111 and 1111111 (in unary representation) for your Turing Machine.

(End of Examination Paper)